

Stacks

Chapter 6

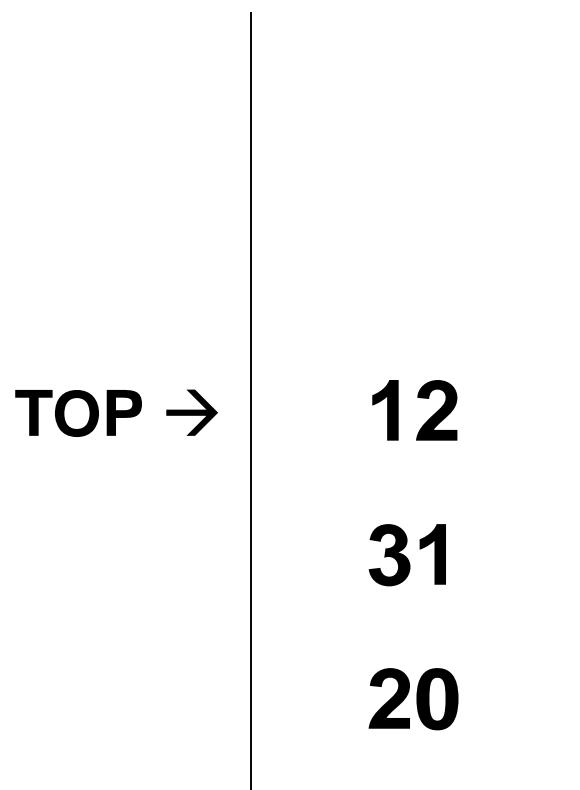
Fundamentals

- A stack is a sequence of data elements (of the same type) arranged one after another conceptually.
- An element can be added to the top of the stack only. (“PUSH”)
- An element can be removed from the top of the stack only. (“POP”)

Applications of Stacks

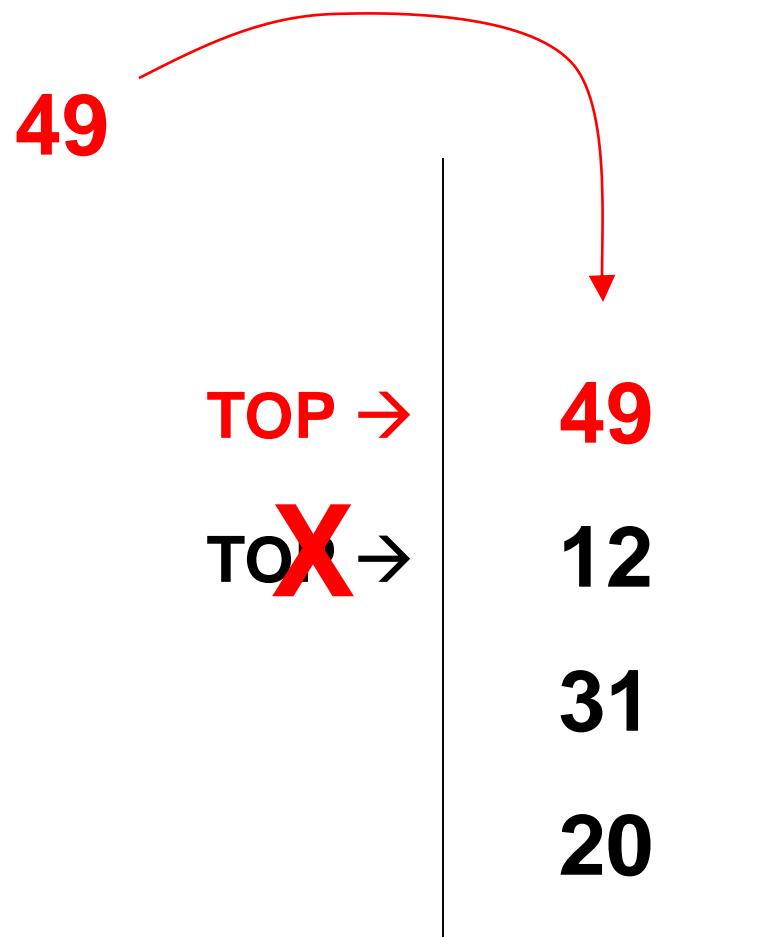
- Operating systems
 - keeping track of method calls in a running program
- Compilers
 - conversion of arithmetic expressions to machine code

Conceptual Picture

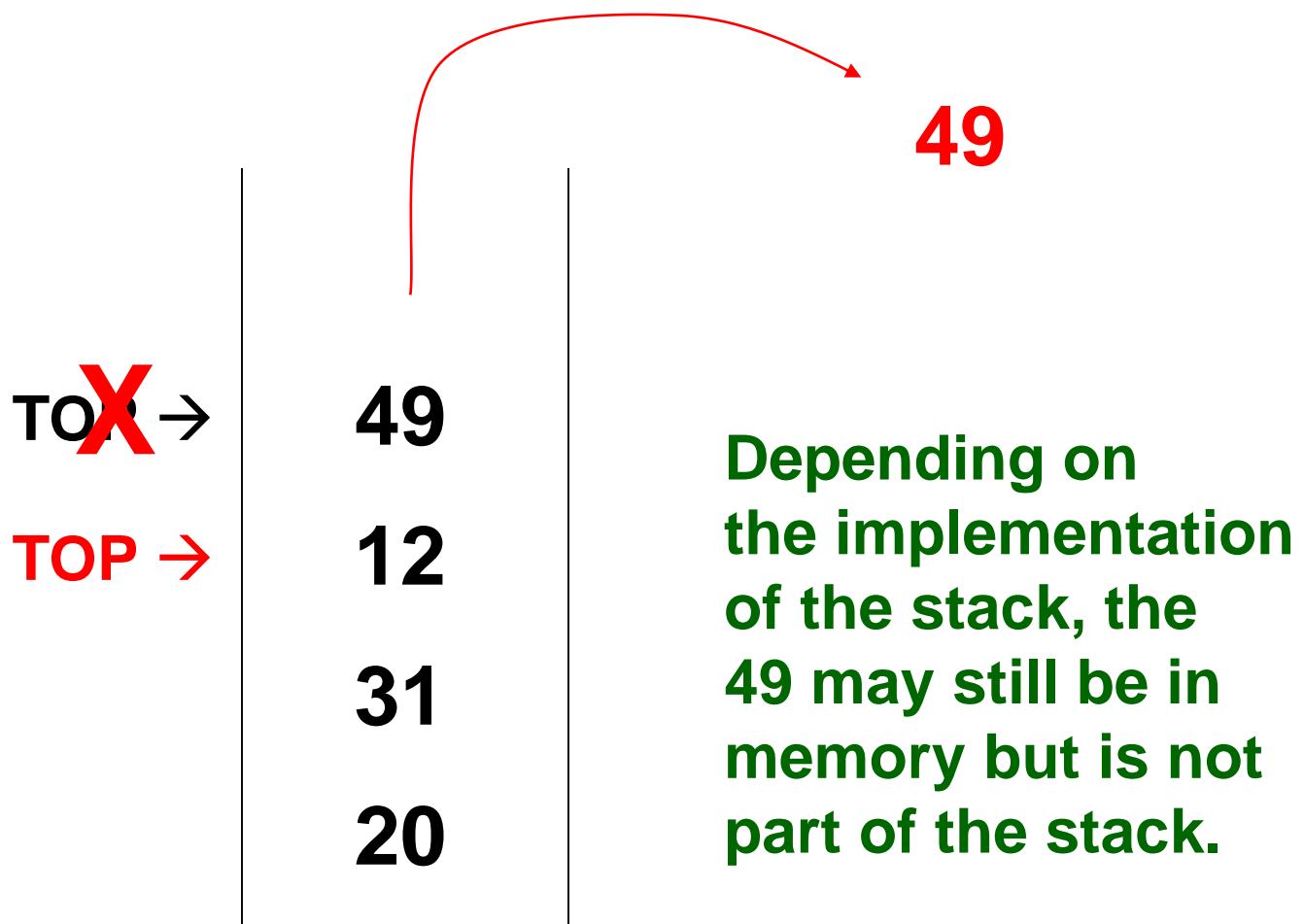


Depending on
the implementation
of a stack, it may
or may not have
a maximum
capacity.

PUSH



POP



Implementation of a Stack

- ARRAYS

12	31	20				
----	----	----	--	--	--	--

top

20	31	12				
----	----	----	--	--	--	--

top

Which is more efficient?

Basic Stack Operations

- **Constructor** – create an empty stack
- **isEmpty** – is the stack empty?
- **push** – push an element on to the top of the stack (if the stack is not full)
- **pop** – remove the top element from the stack (if the stack is not empty)

Extended Stack Operations

- **peek** – examine the top element of the stack without removing it (if the stack is not empty)

```
if not isEmpty()
    temp ← pop()
    push(temp)
    return (temp)
```
- **size** – return the number of elements on the stack
How would you implement this using the basic operations?

An IntStack using Arrays

```
public class IntStack implements  
Cloneable {  
  
    public final int CAPACITY = 100;  
    private int[] data;  
    private int top;  
  
    // IntStack methods (clone not shown)  
}
```

IntStack (arrays) (cont'd)

```
public IntStack()
{
    top = -1;
    data = new int[CAPACITY];
}

public boolean isEmpty()
{
    return (top == -1);
}
```

IntStack (arrays) (cont'd)

```
public void push(int item)
{
    if (top == CAPACITY-1)
        throw new FullStackException();
    top++;
    data[top] = item;
}
```

IntStack (arrays) (cont'd)

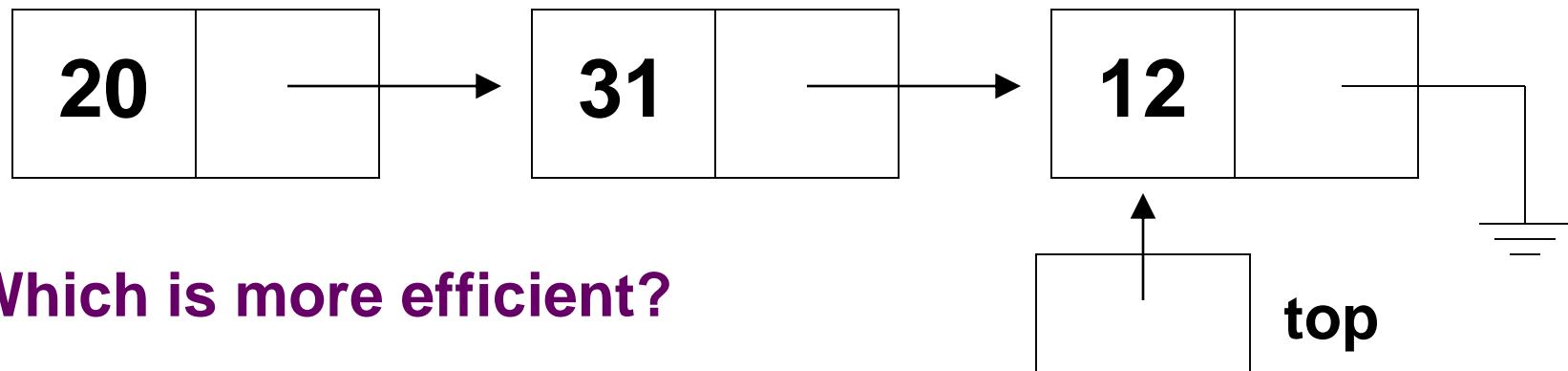
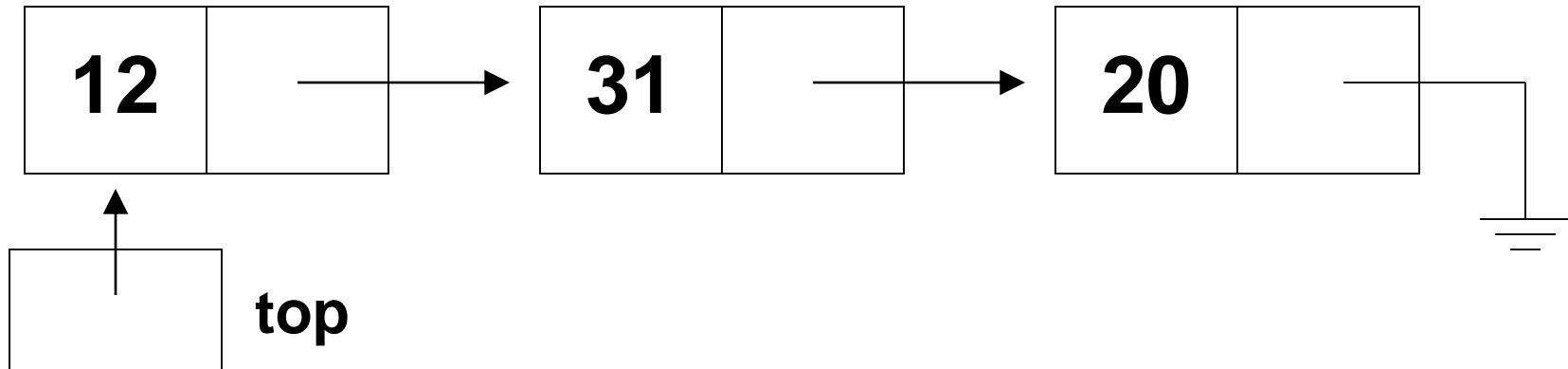
```
public int pop()
{
    int answer;
    if (top == -1) // isEmpty()
        throw new EmptyStackException();
    answer = data[top];
    top--;
    return answer;
}
```

IntStack (arrays) (cont'd)

```
public int peek()
{
    int answer;
    if (top == -1) // isEmpty()
        throw new EmptyStackException();
    answer = data[top];
    return answer;
}
```

Implementation of a Stack

- **LINKED LISTS**



Which is more efficient?

An **IntStack** using Lists

```
public class IntStack implements  
Cloneable {  
  
private IntNode top;  
  
// IntStack methods (clone not shown)  
}
```

IntStack (lists) (cont'd)

```
public IntStack()
{
    top = null;
}
public boolean isEmpty()
{
    return (top == null);
}
```

IntStack (lists) (cont'd)

```
public void push(int item)
{
    IntNode newNode
        = new IntNode(item);
    newNode.setLink(top);
    top = newNode;
}
```

Stack Overflow?

IntStack (lists) (cont'd)

```
public int pop()
{
    int answer;
    if (top == null) // isEmpty()
        throw new EmptyStackException();
    answer = top.getData();
    top = top.getLink();
    return answer;
}
```

IntStack (lists) (cont'd)

```
public int peek()
{
    int answer;
    if (top == null) // isEmpty()
        throw new EmptyStackException();
    answer = top.getData();
    return answer;
}
```

Balanced Parentheses

- An arithmetic expression has balanced parenthesis if and only if:
 - the number of left parentheses of each type is equal to the number of right parentheses of each type
 - each right parenthesis of a given type matches to a left parenthesis of the same type to its left and all parentheses in between are balanced correctly.

Examples

- $(\{A + B\} - C)$
Balanced
- $(\{A + B\} - C\}$
Not balanced
- $(\{A + B\} - [C / D])$
Balanced
- $((\{A + B\} - C) / D))$
Not balanced

Algorithm

CHECK FOR BALANCED PARENTHESES

- Scan the expression from left to right.
 - For each left parenthesis that is found, push on the stack.
 - For each right parenthesis that is found,
 - If the stack is empty, return false
(too many right parentheses)
 - Otherwise, pop the top parenthesis from the stack:
 - If the left and right parentheses are of the same type, discard.
 - Otherwise, return false.

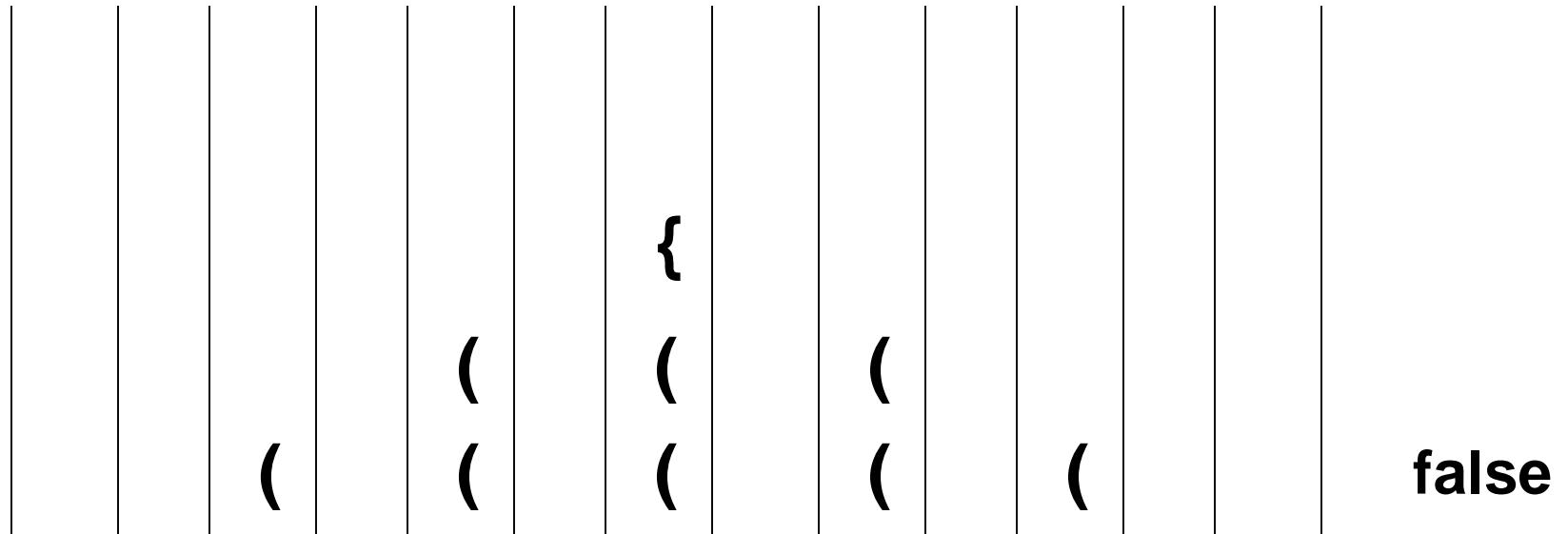
Algorithm (cont'd)

CHECK FOR BALANCED PARENTHESES

- If the stack is empty when the scan is complete, return true.
Otherwise, return false. (too many left parentheses)

Trace

- $((\{A + B\} - C) / D))$
- Stack trace:



Evaluating Expressions

- An expression is fully parenthesized if every operator has a pair of balanced parentheses marking its left and right operands.

- Not fully-parenthesized:

$$3 * (5 + 7) - 9$$
$$(2 - 4) * (5 - 7) + 8$$

- Fully-parenthesized:

$$((3 * (5 + 7)) - 9)$$
$$(((2 - 4) * (5 - 7)) + 8)$$

General Idea

- The first operation to perform is surrounded by the innermost set of balanced parentheses.
- Example: $((3 * (5 + 7)) - 9)$ First op: +
- By reading expression from left to right, first operator comes immediately before first right parenthesis.
- Replace that subexpression with its result and search for next right parenthesis, etc.
- Example: $((3 * 12) - 9) = (36 - 9) = 27$

General Idea (cont'd)

- How do we keep track of operands and operators as we read past them in the expression from left to right?
- Use two stacks:
one for operands and one for operators.
- When we encounter a right parenthesis,
pop off one operator and two operands,
perform the operation, and push the result
back on the operand stack.

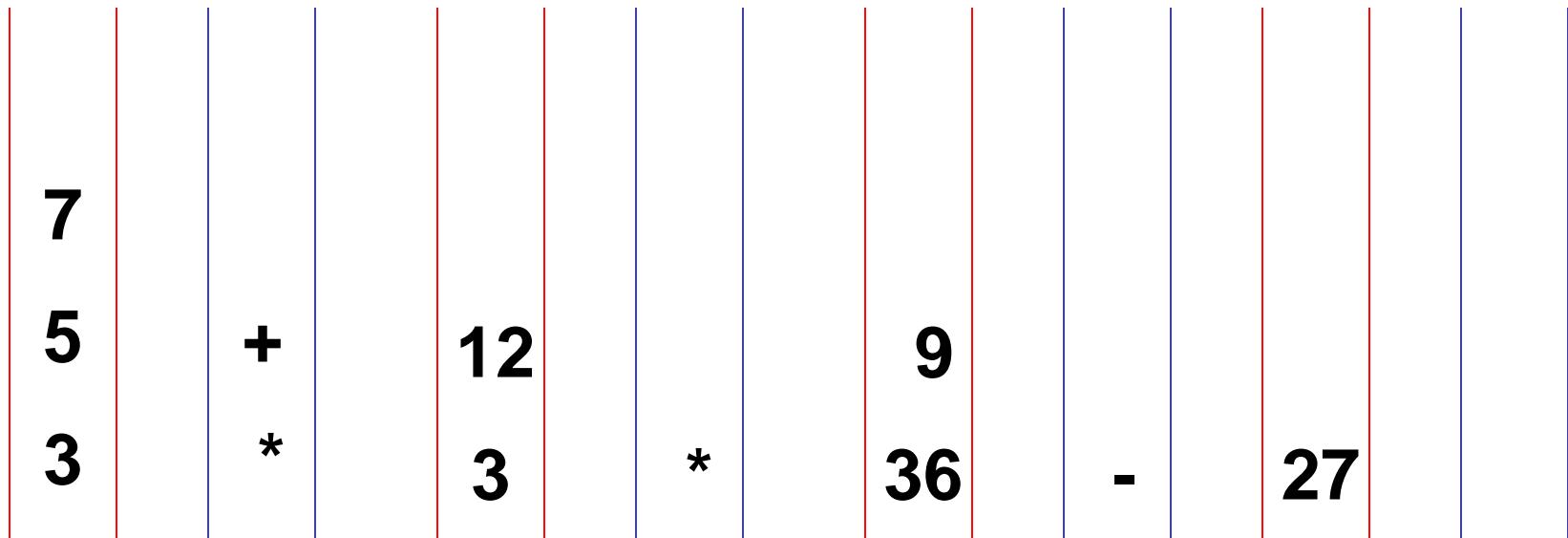
Trace

- $((3 * (5 + 7)) - 9)$

- Stack traces:

operands

operators



ANSWER

Algorithm

- Let each operand or operator or parenthesis symbol be a token.
- Let NumStack store the operands.
- Let OpStack store the operations.
- For each token in the input expression do
 - If token = operand, NumStack.push(token)
 - If token = operator, OpStack.push(token)

Algorithm (cont'd)

If token = “)”,

operand2 \leftarrow NumStack.pop()

operand1 \leftarrow NumStack.pop()

operator \leftarrow OpStack.pop()

result \leftarrow **operand1** **operator** **operand2**

 NumStack.push(**result**)

If token = “(”, ignore token

- After expression is parsed,
answer \leftarrow NumStack.pop()

Arithmetic Expressions

- **Infix notation:**
operator is between its two operands

$3 + 5$ $(5 + 7) * 9$ $5 + (7 * 9)$

- **Prefix notation:**
operator precedes its two operands

$+ 3 5$ $* + 5 7 9$ $+ 5 * 7 9$

- **Postfix notation:**
operator follows its two operands

$3 5 +$ $5 7 + 9 *$ $5 7 9 * +$

NO

PAREN-
THESES

Precedence of Operators

- Multiplication and division (higher precedence) are performed before addition and subtraction (lower precedence)
- Operators in balanced parentheses are performed before operators outside of the balanced parentheses.
- If two operators are of the same precedence, they are evaluated left to right.

Example

- Infix expression:

A + B * (C * D – E / F) / G – H

6 4 1 3 2 5 7

- What is its prefix equivalent?

- + A / * B - * C D / E F G H

- What is its postfix equivalent?

A B C D * E F / - * G / + H -

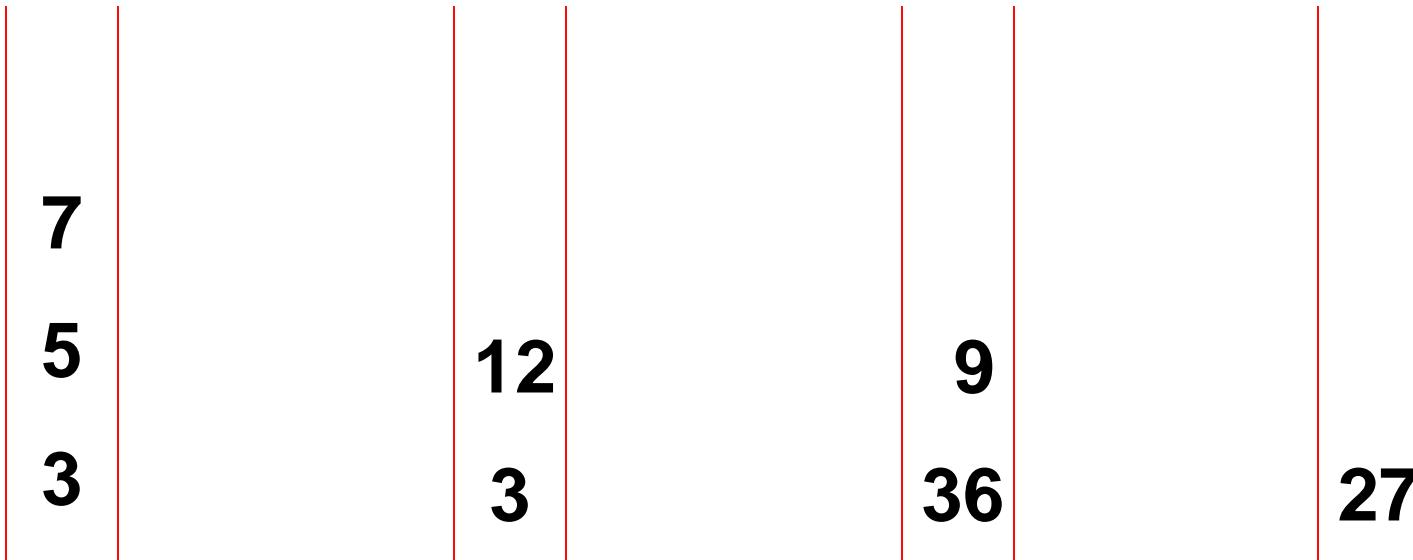
1 2 3 4 5 6 7

Evaluating a Postfix Expression

- Let each operand or operator be a token.
- Let NumStack store the operands.
- For each token in the input expression do
 - If token = operand, push(token)
 - If token = operator,
 - operand $\textcolor{red}{2}$ \leftarrow pop()
 - operand $\textcolor{red}{1}$ \leftarrow pop()
 - result \leftarrow operand $\textcolor{red}{1}$ operator operand $\textcolor{red}{2}$
 - push(result)
- answer \leftarrow pop()

Trace

- **INFIX:** $3 * (5 + 7) - 9$
- **POSTFIX:** $3\ 5\ 7\ +\ *\ 9\ -$
- **Stack traces:** **operands**



ANSWER

Translating Infix to Postfix

Fully-Parenthesized Expressions

- Let each operand, operator, or parenthesis be a token.
- Let OpStack store the operators.
- Let postfix string P = "" (empty string)
- For each token in the input expression do
 - If token = operand, append operand to P
 - If token = operator, push(token)
 - If token = ")", append pop() to P
 - If token = "(", ignore

Trace

- Infix: $((3 * (5 + 7)) - 9)$

Stack (sideways)

empty

*

*

* +

* +

*

empty

-

-

empty

Postfix String

3

3

3 5

3 5

3 5 7

3 5 7 +

3 5 7 + *

3 5 7 + *

3 5 7 + * 9

3 5 7 + * 9 -

Another Example

Infix: $((3 * (5 + 7)) - 9)$

Postfix: 3 5 7 + * 9 -

Infix: $(((2 - 4)^*(5 - 7)) + 8)$

Postfix: 2 4 - 5 7 - * 8 +

Another Example

Infix: $((3 * (5 + 7)) - 9)$

2 1 3

Postfix: 3 5 7 + * 9 -

1 2 3

Infix: $(((2 - 4)^*(5 - 7)) + 8)$

1 3 2 4

Postfix: 2 4 - 5 7 - * 8 +

1 2 3 4

Translating Infix to Postfix

General Expressions

- Define a precedence function
- prec : token $\rightarrow \{0,1,2,3\}$
- let “\$” represent empty stack

token	precedence
“\$”	0
“(”	1
“+”, “-”	2
“*”, “/”	3

Translating Infix to Postfix

General Expressions

- Let each operand, operator, or parenthesis be a token.
- Let OpStack be a character stack that stores the operators or other special symbols (“(” and “\$”).
- Let postfix string P = “” (empty string)

Translating Infix to Postfix (cont'd)

General Expressions

- 1. push("\$")**
- 2. For each token in the input expression do**
 - a. If token = operand,
append token to P**
 - b. if token = "(",
push(token)**

Translating Infix to Postfix (cont'd)

General Expressions

c. if **token = ")"**,

topOp ← pop()

while topOp ≠ "("

append topOp to P

topOp ← pop()

Translating Infix to Postfix (cont'd)

General Expressions

d. if **token = operator**,

topOp ← peek()

while prec(topOp) ≥ prec(token)

 append pop() to P

topOp ← peek()

push(token)

Translating Infix to Postfix (cont'd)

General Expressions

3. At end of infix expression,

topOp ← pop()

while topOp ≠ “\$” do

append topOp to P

topOp ← pop()

Trace

A + B * (C * D - E / F) / G - H

- **Stack (sideways)** **Postfix String**

\$

\$

\$ +

\$ +

\$ + *

\$ + * (

\$ + * (

A

A

A B

A B

A B

A B C

Trace (cont'd)

A + B * (C * D - E / F) / G - H

- **Stack (sideways) Postfix String**

\$ + * (A B C

\$ + * (* A B C

\$ + * (* A B C D

\$ + * (- A B C D *

\$ + * (- A B C D * E

\$ + * (- / A B C D * E

\$ + * (- / A B C D * E F

Trace (cont'd)

A + B * (C * D - E / F) / G - H

- Stack (*sideways*)

\$ + * (- /

\$ + *

\$ + /

\$ + /

\$ -

\$ -

empty

- Postfix String

A B C D * E F

A B C D * E F / -

A B C D * E F / - *

A B C D * E F / - * G

A B C D * E F / - * G / +

A B C D * E F / - * G / + H

A B C D * E F / - * G / + H -